### **Chapter 03: Computer Arithmetic**

Lesson 04: Arithmetic Operations— Multiplication of Integer numbers

# Objective

- Understand the Computer arithmetic operations in Multiplication with unsigned numbers
- Time Taken in Multiplication
- Signed Operand Multiplication of Integers

#### **Multiplication Process**

## **Multiplication**

- Unsigned integer multiplication—handled in a similar manner to the way we multiply multidigit decimal numbers
- The first input to the multiplication
  is multiplied by each bit of the second input
  separately, and the results added by
  binary multiplication

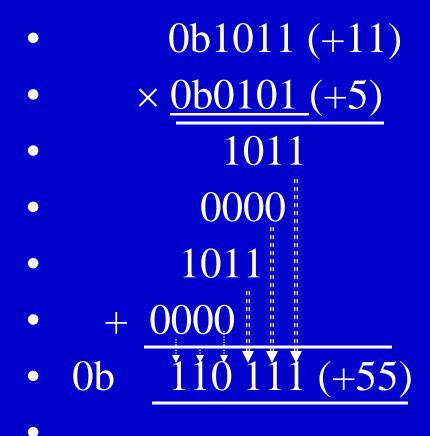
### **Multiplication Process**

- Simplified by the fact that the result of multiplying a number by a bit is either the original number or 0
- Hardware less complex

### **Multiplication Process 11 × 5**

- Multiplying 11 (0b1011), multiplicand (Y) by 5 (0b0l01), multiplier (X)
- First, 0b1011 is multiplied by each bit of 0b0l01 to get the partial products shown
- Then, the partial products are added to get the final result

### **Example of Multiplication Process**



## **Multiplication Process**

#### • Note that each

'Successive partial product is shifted one position to the left to account for the differing place values of the bits in the second input'

#### **Multiplication Circuit**

# **Example** of a Multiplication Circuit

- A method for implementation of the product of two 8-bit numbers using a sequential circuit and one number 8-bit adder
- Assume that two 8-bit registers, A

   (accumulator) and M (multiplier) are used
   for addition
- *A-M* 16-bit combination of two registers for the partial product at each step 0 to Step 7

## **Example of a Multiplication Circuit**

- Let Y (multiplicand) = 0b 10111011 (an unsigned number 187 decimal)
- Let X (multiplier) = 0b 01010101 (an unsigned number 85 decimal)

## **Step 1 of a Multiplication Process**

- 1. The addition (denoted by step A) is done 8 times
- Shift (denoted by step B) is also done 8 times

## **Step 2 of a Multiplication Process**

- Shift is to the right when we use multiplier bits from msb down to lsb during steps 0 to 7
- Shift takes 17-bits into account: the carry plus A-bits and M-bits
- C will shift to msb of A in step B

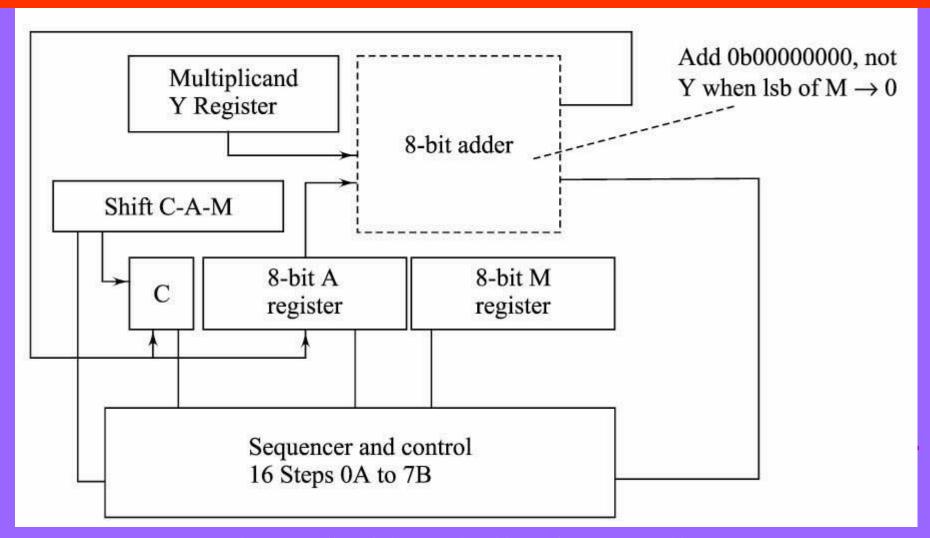
## **Step 3 of a Multiplication Process**

3. Each step has two parts, *A* and *B*, and a C-flag of 1-bit stores the -carry out shift-right from maximum significant bit (msb) at *A* 

### **Step 4 of a Multiplication Process**

- 4. Two registers shift and one addition per step = 16 shifts of 8-bit registers and 8 additions of 8-bit registers = 24 operations + 4 clear plus load M = 28
  - In the present example, the total number of expected operations = 24, 4 addition operations did not occur due to the nature of the multiplier

#### Sequential circuit and 8-bit multiplier implementation by 16 steps (8-cycles) of addition and shifts



# Steps in two Dimesnsional Array to get partial product

Step	C-flag *	First Register for A	Second Register for M	Action Taken	Number of operations (instructions)	
Start	0 0b0000000		0b000000 <u>0</u>		3 for clearing $C$ , A and $M$	
	0	060000000	0b0101010 <u>1</u>	Load Multiplier X in M	1	
Step 0A	A 0 10111011		0101010 <u>1</u>	Add multiplicand $Y$ in $A$ , result in C- $A$	1	
Step 0B	0 0B 0 01011101		10101010	Shift C-A-M	2	
Step 1A	A 0 01011101 10101010		10101010	Do not add (lsb of $M = 0$ )	0 (1)	
Step 1B	0 1B 0 00 101110		1101010 <u>1</u>	Shift C-A-M	2	
Step 2A	p 2A 0 11101001		11010101	Add Y, result in C-A	1	
Step 2B	2B 0 01110100 1110101		→ <u>11101010</u>	Shift	2	

# Steps in two Dimesnsional Array to get partial product

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Step 3A	0 ~	01110100	1110101 <u>0</u>	Do not add $(lsb = 0)$	0(1)	
Step 3B	0	<b>→</b> 00111010	∽ <u>01110101</u>	Shift C-A-M	2	
Step 4A	0 /	11110101	01110101	Add Y, result in C-A	1	
Step 4B	0	01111010 10111010		Shift	2	
Step 5A	0	01111010	10111010	Do not add $(lsb = 0)$	0(1)	
Step 5B	0	00111101 01011101		Shift C-A-M	2	
Step 6A	0	11111000	01011101	Add Y, result in C-A	1	
Step 6 B	0	01111100	→ 00101110	Shift C-A-M	2	
Step 7 A	0 /	01111100	0010111 <u>0</u>	Do not add $(lsb = 0)$	0(1)	
Step 7 B	0	00111110	→ 00010111	Shift C-A-M	2	
Answer	0	0011 1110 0001 0111= 0x3e17		Decimal 15895	Total 24 (28)	

\* after the shift-left from the msb of A.

# Two dimensional array of Full adders to get partial products

Multiplier X Bits	Input Y to Full Adders	Cycle Number	7	6	5	4	3	2	1	0
200000 - 2000 - 3000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000 - 2000	0000		0	0	0	0	0	0	0	0
1 (lsb)	00001011	0					1	0	1	1
0 0000		1				0	0	0	0 🔫	
1	000101100	2			1	0	1	1 -		8-1 
0 (msb)	0000	3		0	0	0	0 -			
			FA	FA	FA	FA	FA	FA	FA	FA
			FA	FA	FA	FA	FA	FA	FA	
			FA	FA	FA	FA	FA	FA		
			FA	FA	FA	FA	FA			
			FA	FA	FA	FA				
			FA	FA	FA					
			FA	FA						
			FA							
										0

Arrow in a column shows implementation of shift.

#### **Time taken in Multiplication**

# **Example: Compute the time taken for an 8bit multiplication using the circuit**

- Assume that an 8-bit adder adds in 0.008 µs, and a shift of carry bit + 8-bit accumulator and multiplier registers' shift by one bit after each addition takes 0.002 µs.
- Assume that the time for other operations is 0.001µs

## **Solution**

- The multiplication circuit design has 8-bit addition and shift lefts, both 8-times
- There will be 8-bit additions in 8-bit multiplier and 8-times shifts

## **Solution**

- Time in 2 registers and carry clear=  $0.003 \ \mu s$
- Time in register loads  $= 0.001 \ \mu s$
- Time taken for 8-bit addition before the shift =  $0.008 \ \mu s.$
- Time taken for 1-bit shift  $= 0.002 \ \mu s$
- Time taken for 8-add and 8-shifts =  $0.003 \ \mu s + 0.001 \ \mu s + 8 \times (0.008 + 0.002) \ \mu s$  $= 8 \times 10 \ ns + 4 = 84 \ ns$
- Multiplication Time = 84 ns

### **Signed operand multiplication of integers**

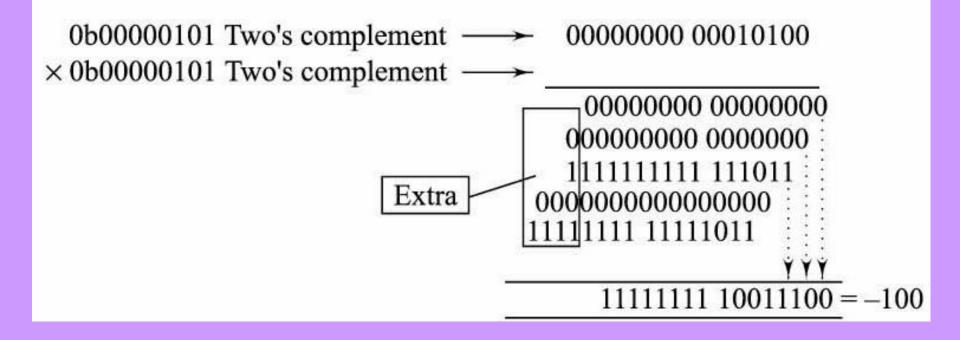
## **Signed Integer Multiplication**

- Signed integer multiplication handled in a manner similar to the way unsigned integers
- Multiply multidigit decimal numbers and accumulate the partial products

### **Multiplication Process**

 However, for an n-bit signed multiplier, there should be a sign extension up to 2<sup>n</sup> bits and we must find the two's complement of both

## **Multiplication of signed numbers**



## Summary

### We learnt

- Multiplication process is a process in which the first input to the multiplication is multiplied by each bit of the second input separately
- The results added by binary multiplication
- Successive partial product shifted one position to the left to account for the differing place values of the bits in the second input

End of Lesson 4 on Arithmetic Operations— Multiplication of Integer numbers