Chapter 02: Computer Organization

Lesson 07 Operating Systems

Objective

- Understand concepts of OS, Multiprogramming, Protection and virtual memory
- Learn OS control of the resources and time slice
- Understand Privilege and user modes

Operating System

The operating system

 Simply another program, one that knows about all the hardware in the computer, with one exception — it runs in privileged (or supervisor) mode

Privileged Mode

 Allows it access to physical resources that user programs (tasks) cannot control and gives it the ability to start or stop the execution of user programs

OS Responsibilities

- Managing the physical resources of the system,
- Loading and executing programs
- Managing devices and network and
- Controlling through *device drivers* the I/O devices

Multiprogrammed operating system

• It presents the illusion that multiple programs are running simultaneously by switching between programs very rapidly

Multiprogramming

Multiprogramming

 Most computer systems support multiprogramming (also called multitasking)— a technique that allows the system to present the illusion that multiple programs are running on the computer simultaneously, even though the system may only have one processor

Multi programmed System

• User programs do not need to know which other programs are running on the system at the same time they are, or even how many other programs there are

Multiuser computers

• Many multiprogrammed computer are also *multiuser* and allow more than one user to be logged in to the computer at once

Multiuser computers

 Multiuser systems require that the operating system not only protect programs from accessing each other's data but prevent users from accessing data that is private to other users

Multiprogrammed operating system

• Each program allowed to execute for a fixed amount of time, known as a *timeslice*

Context Switch

- When a program's timeslice ends, the operating system stops it, removes it from the processor, and loads another program into the processor
- This process known as a *context switch*

Context switch

• The operating system copies the contents of the currently running program's register file (sometimes called the program's context) into memory, and then copies the contents of the next program's register file out of memory and into the register file

Context Switch

 Programs cannot tell that a context switch has been performed—to them, it looks like they have been continuously running on the processor

Context-switch Period

- 60 times per second, making timeslices 1/60th of a second,
- Lately systems have started to context-switch more frequently

Fast Context Switch Time

• This has caused problems with programs that expect timeslices to be 1/60th of a second and use that information to time events or determine performance

Example 1

• By context-switching 60 or more times per second, a computer can give each program an opportunity to execute sufficiently frequently that the system can prevent the illusion that a moderate number of programs are executing simultaneously on the system

Example 2

 Obviously, as the number of programs on the system increases, this illusion breaks down—if the system is executing 120 programs, each program may only get a timeslice once every 2 seconds, which is enough of a delay for us to notice.

Increase of the running time in Multiprogramming

- Running time of applications increase
- The resources of the system are shared among all of the programs running on it

Protection

Main requirements of a multiprogrammed operating system

• One is that OS must provide *protection* between programs running on the computer

Main requirements of a multiprogrammed operating system

 Result of any program running on a multiprogrammed computer must be the same as if the program was the only program running on the computer

Operating system

 OS and hardware provide *protection* for programs, preventing any program from accessing another program's data unless the two programs have specifically arranged to access each other's data

Protection

- Programs must not be able to access other programs' data and must be confident that their data will not be modified by other programs.
- Similarly, programs must not be able to interfere with each other's use of the I/O subsystem

Use of Virtual memory

 One technique that operating systems use to protect each program's data from other programs is *virtual memory*

Use of Virtual memory

 Allows each program to operate as if it were the only program running on the computer by *translating* memory addresses that the program references into the addresses used by the memory system

Virtual memory system

- Ensures that two programs' addresses don't translate into the same address
- Programs can be written as if they were the only program running on the machine, since no program's memory references will access data from another program

Control of physical resources

Control of physical resources

 Providing protection in a multiprogrammed or multiuser system requires that the operating system control the physical resources of the computer, including the processor, the memory, and the I/O devices

Without Protection

 User programs could access any of the memory or other storage on the computer, gaining access to data that belongs to other programs or other users

With Protection

• Allows the operating system to prevent more than one program from accessing an I/O device, such as a printer, at one time

Privileged mode

Privileged mode

- Ensures that the operating system is the only program that can control the system physical resources,
- User programs execute in *user mode* (sometimes called unprivileged mode)

Context switches and memory allocations

- Certain tasks, such as accessing an I/O device, require that a program be in privileged mode
- If a user-mode program tries to perform one of these tasks, the hardware prevents it from doing so and signals an error

User-mode programs

• If want to do something that requires privilege mode, they send a request to the operating system, known as a *system call*, which asks the operating system to do the operation for them

User System calls

- If the operation is something that the user program is allowed to do, the operating system performs the operation and returns the result to the user
- Otherwise, signals an error

User Interface

 Because OS controls the physical resources of the computer, it is also responsible for the low-level user interface

Example 1

• When a user presses a key or otherwise sends input to the computer, the operating system is responsible for determining which program should receive the input and sending the input value to that program

Example 2

• When a program wants to display some information for the use such as printing a character on the monitor, it executes a system call to request that the operating system display the data.

Summary

We learnt

- OS control of resources,
- Multiprogramming,
- Context Switch
- Protection,
- Time slice
- Privilege Mode and user mode

End of Lesson 7 on **Operating Systems**